



# On the Origin of Trust: Struggle for Secure Cryptography

Anne Canteaut

## ► To cite this version:

Anne Canteaut. On the Origin of Trust: Struggle for Secure Cryptography. Dot Security 2016, Apr 2016, Paris, France. hal-01401311

**HAL Id: hal-01401311**

**<https://inria.hal.science/hal-01401311>**

Submitted on 23 Nov 2016

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# On the Origin of Trust: Struggle for Secure Cryptography

Anne Canteaut

Inria, project-team SECRET, Paris, France

<http://www.rocq.inria.fr/secret/Anne.Canteaut/>

# Attacks

**In most attacks, cryptography is bypassed.**

“I am not aware of any major world-class security system employing cryptography in which the hackers penetrated the system by actually going through the cryptanalysis.” [Adi Shamir 2002]

# Attacks

**In most attacks, cryptography is bypassed.**

“I am not aware of any major world-class security system employing cryptography in which the hackers penetrated the system by actually going through the cryptanalysis.” [Adi Shamir 2002]

“I do not need a trophy to tell myself that I am the best.”  
[Zlatan Ibrahimovic 2013]

# Disasters...





[Home](#) [India](#) [World](#) [Business](#) [Technology](#) [Sports](#) [Entertainment](#) [Lifestyle](#) [Ec](#)

**TRENDING#** [Maharashtra Drought](#) [IPL 2016](#) [Assam Elections 2016](#) [West Bengal electio](#)

[Home](#) >> [Sport](#)



40

0

1

41



## Champions League: Ibrahimovic misses spot kick, PSG pay the penalty against Man City

*Fernandinho blasts in Manchester City's equalizer to stun the PSG supporters (Reuters)*

Thu, 7 Apr 2016-10:01am , Paris , Reuters

# TLS/SSL attacks

- biases in RC4 [AlFardam et al. 13]
- Logjam [Adrian et al. 15]: weak Diffie-Helman
- Sloth [Bhargavan, Leurent 16]: collisions in MD5

# Attack against MIFARE

**COMPUTERWORLDUK**  
THE VOICE OF BUSINESS TECHNOLOGY

FeaturesTechnologyIT BusinessCareersTut

Resources



Home > News > IT Vendors > Questions raised about Oyster card security

## Questions raised about Oyster card security

Its RFID chip is cracked by researchers

Network World and  
Computerworld UK staff  
March 7, 2008

re Smartcards with encrypted RFID chips, including London's Oyster fare card, might not be as secure as previously thought.

**Can we trust cryptographers?**



# Rule #1

**When cryptographers claim that a primitive is broken,  
don't use it.**

# Rule #1

**When cryptographers claim that a primitive is broken,  
don't use it.**

**But those guys are paranoid!**

# Rule #1

**When cryptographers claim that a primitive is broken,  
don't use it.**

**But those guys are paranoid!** True.

Cryptanalysis of the full Spritz [Banik, Isobe 16]:

“We need approximatively  $2^{1247}$  assignments to recover the internal state.”

# Rule #1

**When cryptographers claim that a primitive is broken,  
don't use it.**

**But those guys are paranoid!** True.

Cryptanalysis of the full Spritz [Banik, Isobe 16]:

“We need approximatively  $2^{1247}$  assignments to recover the internal state.”

$$2^{1247} \geq (\# \text{ atoms in the universe} )^4$$

# Broken?

A good primitive must behave as a function chosen at random from the set of all functions (with the same characteristics).

# Broken?

A good primitive must behave as a function chosen at random from the set of all functions (with the same characteristics).

Spritz [Rivest, Schulz 15]:

Spritz generates a pseudo-random sequence from a secret state, chosen out of  $2^{1730}$  possibilities.

# Broken?

A good primitive must behave as a function chosen at random from the set of all functions (with the same characteristics).

Spritz [Rivest, Schulz 15]:

Spritz generates a pseudo-random sequence from a secret state, chosen out of  $2^{1730}$  possibilities.

**Attack:** the internal state can be recovered with  $2^{1247}$  trials  
→ much better than brute-force

# Hash functions

$$H : \{0, 1\}^* \longrightarrow \{0, 1\}^n$$

## Second preimage:

Given  $m$ , find a message  $m'$  such that  $H(m') = H(m)$ .

Generic algorithm: Try  $2^n$  random messages.



# Hash functions

$$H : \{0, 1\}^* \longrightarrow \{0, 1\}^n$$

## Second preimage:

Given  $m$ , find a message  $m'$  such that  $H(m') = H(m)$ .

Generic algorithm: Try  $2^n$  random messages.

## Collision:

Find two messages  $m$  and  $m'$  such that  $H(m) = H(m')$ .

Generic algorithm: Select  $2^{n/2}$  random messages.

**But this is not relevant in our applications...**

Finding collisions is not an issue in key-exchange protocols.

**But this is not relevant in our applications...**

Finding collisions is not an issue in key-exchange protocols.

Sloth attack against TLS [Bhargavan, Leurent 16]:  
exploits collisions in MD5!

# But these attacks are not practical...

- Attacks reveal unexpected weaknesses.

## But these attacks are not practical...

- Attacks reveal unexpected weaknesses.
- Attacks always get better; they never get worse.

## But these attacks are not practical...

- Attacks reveal unexpected weaknesses.
- Attacks always get better; they never get worse.

**If cryptographers say that it is broken, don't use it.**

**What if they don't say  
that it is broken?**

# Is there any difference between

- **AES** (NIST FIPS 197)
- **Crypto-1** (MIFARE Classic encryption)
- **Dual-EC-DRBG** (NIST SP 800-90A)



# Is there any difference between

- **AES** (NIST FIPS 197)
- **Crypto-1** (MIFARE Classic encryption)
- **Dual-EC-DRBG** (NIST SP 800-90A)

**AES has been standardized after an open competition  
(1997-2001)**

# Hash function competition (SHA-3)

Oct 2008 submission deadline

→ 64 candidates received by the NIST

Dec 2008 51 candidates in the 1st round

Feb 2009 1st SHA-3 conference

# Let's start the struggle!

Abacus	Neil Sholer	in round 1	2nd-preimage	
ARIRANG	Jongin Lim	in round 1		
AURORA	Masahiro Fujita	in round 1	2nd preimage	
Blender	Colin Bradbury	in round 1	collision, preimage	near-collision
Boole	Greg Rose	in round 1	collision	
Cheetah	Dmitry Khovratovich	in round 1		length- extension
CHI	Phillip Hawkes	in round 1		
CRUNCH	Jacques Patarin	in round 1		length- extension
DCH	David A. Wilson	in round 1	collision	
Dynamic SHA	Xu Zijie	in round 1	collision	length- extension

[http://ehash.iaik.tugraz.at/wiki/The\\_SHA-3\\_Zoo](http://ehash.iaik.tugraz.at/wiki/The_SHA-3_Zoo)

# Hash function competition (SHA-3)

- Oct 2008 submission deadline
  - 64 candidates received by the NIST
- Dec 2008 51 candidates in the 1st round
- Feb 2009 1st SHA-3 conference
- July 2009 14 candidates in the 2nd round
- Aug 2010 2nd SHA-3 conference
- Dec 2010 5 finalists
- Mar 2012 3rd SHA-3 conference
- Oct 2012 winner announced (Keccak)

# Prize for the best cryptanalysis

## Third cryptanalysis prize

---

30 September 2009

We announce the third prize for the most interesting cryptanalysis of KECCAK. The results must be publicly available on an URL that is sent to keccak -at- noekeon -dot- org **before December 5, 2009** at 23:59 GMT+1 (i.e., before Sinterklaas or Saint Nicolas).

The third prize consists of beer, like the first one. This time we offer Lambic beers that according to myth can only be brewed in the surroundings of Brussels thanks to wild yeast and mysterious bacteria that would not occur anywhere else. Anyway, the prize is a case with 24 (the new number of rounds in KECCAK-f) bottles of Lambic-based beers from breweries such as Cantillon, Girardin, and 3 Fonteinen.

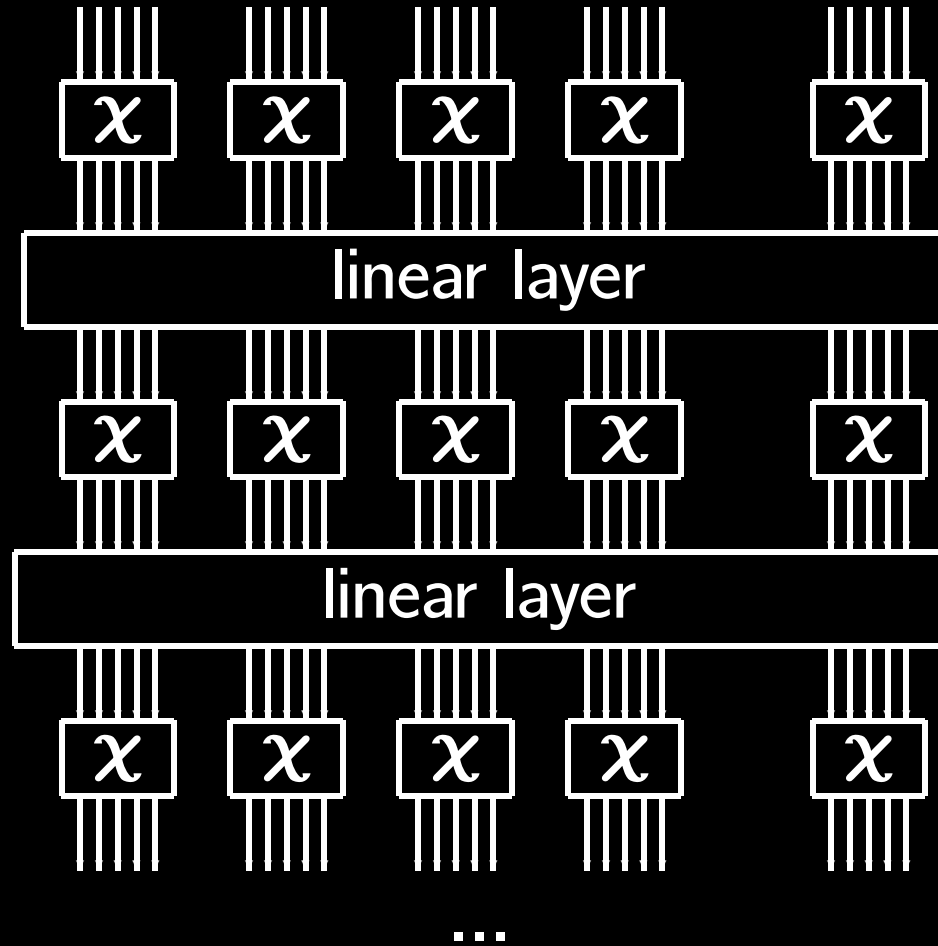
[http://keccak.noekeon.org/third\\_party.html](http://keccak.noekeon.org/third_party.html)

# Prize for the best cryptanalysis

[Boura, Canteaut 2011]: distinguisher on the inner permutation of Keccak with complexity  $2^{1575}$  (instead of  $2^{1600}$ ).



## Round-reduced versions



In Keccak, 24 rounds

# How many rounds can we break?

**SHA-3 (24 rounds):**

collisions up to 5 rounds [Dinur, Dunkelman, Shamir 2013]



# How many rounds can we break?

## AES-128 (10 rounds):

5 rounds	$2^{46}$	Daemen, Rijmen 1998
6 rounds	$2^{71}$	Daemen, Rijmen 1998
6 rounds	$2^{48}$	Ferguson et al. 2000
7 rounds	$\simeq 2^{128}$	Gilbert, Minier 2000
7 rounds	$2^{117}$	Lu, Dunkelman, Keller, Kim 2008
7 rounds	$2^{110}$	Mala et al. 2010
7 rounds	$2^{99}$	Derbez, Fouque, Jean 2013

# Rule #2

**No public analysis, no trust**

## Examples:

- **Crypto-1 (Mifare):** proprietary design
- **Simon, Speck [NSA 2015]:** no design rationale

# Conclusion

**Public analysis** is the only reliable security argument